# Approximate Implicitization and CADtype intersection algorithms 

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# The starting point is: <br> Geometry representation in CAD-systems 

## Standardized in ISO 10303 STEP in the early 1990es.

- Degree 1 and 2 algebraic curves and surfaces + torus
- NonUniform Rational B-Spline (NURBS) curves and surfaces
- Piecewise rational/polynomial curves and surfaces
- Frequently cubic/bi-cubic, but also higher degrees allowed (and in use)
- Volumes represented by description of the outer shell (and inner shell(s))
- A shell is represented by a patchwork of surface pieces
- A shell is not required to be watertight, small tolerances controlled gaps allowed
- A surface patch is limited by edges, the edge is limited by two vertices
- Double precision floating point arithmetic used

The representation is not precise!

## Why is CAD-geometry represented this way?

- ISO 10303 standardize the ideas of the late 1980s.

■ Monolithic 3D applications dominated

- 3D CAD was still immature

■ Computers are now at least 3 orders of magnitude faster

- Memory sizes are now 2 to 3 orders of magnitude larger
- Consequently 3D CAD is far from optimal but

■ It has penetrated all branches of industry

- The industry has invested heavily in CAD
- The CAD-industry has merged into a few dominant vendors

■ Current CAD is good enough for the average user but not for high-end industries such as aerospace, automotive and oil \& gas industries

## CAD design often produce near singular transitions between surfaces

- Patchwork of surface to build a larger smooth surfaces
- Transition between blending surface and mother surfaces

■ Design intent and what the user believes happens: Tangent continuity

- Result: Small gaps and near tangent continuous



## What does near singular mean?

- Seen from far away -

■ Intersection interval

- Zooming by a factor of 10x

■ No intersection

- Many computer displays have less than $1200 \times 1600$ pixels
■ When visualizing an object of size 1000 mm , the smallest visible details will be approximately 1 mm
- Production tolerance in most cases significantly smaller....
- The displayed image also distorted by tessellation


## Partial coincidence



## Traditional CAD-type intersection algorithms focus on non-singular intersections

- An intersection curve between two surfaces is transversal when the normals of the intersecting surfaces are non-parallel along the intersection curve
- Sinha's theorem (1985):
- If two smooth surfaces $S_{1}$ and $S_{2}$ intersect in a common loop then there is a point $P_{1}$ inside the loop in $S_{1}$, and there is a point $P_{2}$ inside the loop in $S_{2}$ such that the normal $N_{1}$ in $P_{1}$ is parallel to the normal $N_{2}$ in $P_{2}$.
- If the normal fields of two surfaces do not overlap, no closed intersection loop is possible, and the intersection is transversal.
- Repeated subdivision of surfaces with overlapping normal fields will, provided the intersections curve is transversal, eventually results in subproblems where Sinha's theorem can be applied. (Loop destruction)
- However, when the intersection is near singular it will take a very long time....


## Singular and near singular intersections



2 points!

2 points?
1 singular point?
An interval?
No point?

- The relative position and orientation of curves and surfaces determines if an intersection is:
■ Transversal
- Near singular (tolerance dependent!)
- Singular


## Recursive subdivision to try to make simpler subproblems



■ Each intersection singled out in a simple subproblem

## Recursive subdivision do not efficiently sort out all singular or near singular situations

■ Difficult to decide if sculptured near parallel curves and surfaces intersect or not


- Deep levels of recursion necessary
- Where to subdivide has to be considered with care


## Most CAD-intersection algorithms have no quality guarantee

- Simplistic algorithms are fast \& often produce the correct result
- Intersect triangulations of the surfaces

■ Lattice evaluation - intersect mesh of curves in each surface with the other surface to possibly generate points on all intersection branches
■ Marching/refinement of identified intersection tracks

- Recursive algorithms slower, sometimes extremely slow

■ More calculations, guarantee for clearly transversal intersections
■ Deep levels of recursion in near singular cases

- For singular intersections traditional recursive intersection algorithms will not work (well)
- Cut off strategies necessary to avoid infinite recursion
- Improved approaches needed


## Improvement of intersection algorithms by combining parametric and algebraic representations

- Improved approaches for separating surfaces
- Simplification of intersection problems to the parameter domain of one of the surfaces
■ Determine that two surfaces only touch along a boundary curve

Most often an algebraic surface approximating the part of the surface addressed suffice.

## Approximate implicitization (Dokken 97)

- In stead of the global correct (high degree) algebraic representation we want to find an algebraic approximation to the curve or surface that is closer than a given tolerance in a defined region of interest.
■ Well behaved numeric method "Approximate Implicitization" have been developed
- Proven numeric well behaved rounding error
- High convergence rates
- Use modified LU-decomposition or Singular Value Decomposition

■ Algebraic degree can be considerably lower than the theoretical exact degree. (For bicubic total degree 4 or 6, opposed to the exact degree 18)
■ Sufficiently efficient to be an efficient tool for determining intersection, near intersection or separation of surfaces intersected

- The method is an exact implicitization method if proper algebraic degree chosen (and exact arithmetic used)


## The approximate implicitization factorization

Assume that the surface $\mathbf{p}(s, t)$ has bi-degree $\left(n_{1}, n_{2}\right)$

- Assume that $q$ has total degree $m$ and that $\mathbf{b}$ is a vector containing the unknown coefficients of $q$
- The combination $q(\mathbf{p}(s, t))$ is a polynomial of bi-degree $\left(m n_{1}, m n_{2}\right)$
- Collect basis functions of bi-degree $\left(m n_{1}, m n_{2}\right)$ in $\alpha(s, t)$
- Then $q(\mathbf{p}(s, t))$ can be factorized

$$
q(\mathbf{p}(s, t))=(\mathbf{D b})^{T} \boldsymbol{\alpha}(s, t)
$$

## The factorization

$$
q(\mathbf{p}(s, t))=(\mathbf{D b})^{T} \boldsymbol{\alpha}(s, t) .
$$

- An element in $\mathbf{D}$ is the product of a maximum of $m$ coefficients of $\mathbf{p}(s, t)$ and a constant, where $m$ is the total degree of $q$.
- If $\mathbf{p}(s, t)$ is a Bezier surface of bi-degree $\left(n_{1}, n_{2}\right)$ then $\alpha(s, t)$ is a Bernstein basis of bi-degree $\left(m n_{1}, m n_{2}\right)$.


## Properties of the factorization

$$
q(\mathbf{p}(s, t))=(\mathbf{D b})^{T} \boldsymbol{\alpha}(s, t)
$$

- If $\mathbf{D b}=\mathbf{0}$ and $\mathbf{b} \neq \mathbf{0}$ then $\mathbf{b}$ contains the coefficients of an exact algebraic representation of total degree $m$ of $\mathbf{p}(s, t)$.
- If $\alpha(s, t)$ is a Bernstein basis then $\|\boldsymbol{\alpha}(s, t)\|_{2} \leq 1$, and

$$
|q(\mathbf{p}(s, t))|=\left|(\mathbf{D b})^{\mathrm{T}} \boldsymbol{\alpha}(s, t)\right| \leq\|\mathbf{D} \mathbf{b}\|_{2} .
$$

■ Let $\sigma_{\text {min }}$ be the smallest singular value of $\mathbf{D}$, then

$$
\min _{\|\mathbf{b}\|_{2}=1} \max _{(s, t) \in \Omega}|q(\mathbf{p}(s, t))| \leq \sigma_{\text {min }} .
$$

- Singular value decomposition of $\mathbf{D}$ can be used to find approximate solutions


## The algebraic/parametric combination used for separation of surfaces

Let $\mathbf{p}(s, t),(s, t) \in \Omega_{1}$ and $\mathbf{r}(u, v),(u, v) \in \Omega_{2}$, be two rational surfaces

- Decide that two surfaces do not intersect by finding an algebraic surface $q(x, y, z)=0$ separating the surfaces $q(\mathbf{p}(s, t))>c$ and $q(\mathbf{r}(u, v))<c$.
- Find the approximate algebraic surface by
approximate implicitization



## The algebraic/parametric combination for determining the topology of an intersection

- The intersection of two parametric curves $\mathbf{p}_{1}(\mathrm{~s})$ and $\mathbf{p}_{2}(\mathrm{t})$, can be simplified if implicit representations of at least one of curves exist: $q_{1}(x, y)=0$ and $\mathrm{q}_{2}(\mathrm{x}, \mathrm{y})=0$.
- The combination $q_{1}\left(p_{2}(t)\right)=0$ transforms the intersection of two parametric curves to finding the zeroes of an univariate polynomial
- Easily extended to surfaces use approximate implicitization



## Special use of approximate implicitization in self-intersection

■ As part of a recursive surface self-intersection algorithm, adjacent surface subpatches have to be intersected
■ There will always be an intersection along the common edge
■ An approximate implicit surface following the normal of the surface (or a fixed direction) along the edge between the subpatches is made, and used for deciding if the edge intersection is the only intersection between the subpatches

## An example where traditional recursive intersection algorithms work well

- Intersection of a plane parametric surface and a varying parametric surface producing many intersection loops



## Singular intersection curves and loops



## Surface self-intersection can give complex intersection topology



Self-intersecting bi-cubic B-spline surfaces.

- $39 \times 70$ polynomial pieces
- Single knots

Courtesy think3


Wire frame of surface with self-intersection curves. The self-intersection curves displayed alone to the right

## Parameter domain self-intersection trace



Singular points (vanishing normal)


## More details



## Intersection within intersection



A small self-intersection loop very close to the global self-intersection The distance between the loops is $1 / 10000$ of the width of the parameter domain.


## The GAIA Surface Self-intersection code

- Originally we aimed at two types of self-intersections

■ Global. Two completely different pieces of the surface intersect. Provided that the surface is split into relevant sub surfaces, global self-intersections can be computed as surface-surface intersections..

- Local. A local self-intersection will appear as a small loop or a cusp. The surface normal will become very small in the vicinity of a local selfintersection
- During testing we realized that cusp ridges (curves where the surface normal vanish) are more frequent than expected in self-intersection
■ The ridges do not in general follow constant parameter lines
- Offset surfaces, duct type surfaces
- The self-intersection code uses the GAIA surface surface intersection code and has posted new challenges to this code


## Future work

- Improve the intersection code by:

■ Further testing and debugging

- Improve speed of approximate implicitization

■ Implement new strategies for the combination of recursive subdivision and approximate implicitization

- GPU-acceleation
- To better understand what is going on we will in our Paralle3D project (www.sintef.no/parallel3d)
■ Improve visualization tools to be able to zoom into more detail. The current viewers do not allow fine enough tessellation
■ Combine viewers for the parameter domain and 3D
- Combine viewers for algebraic and parametric surfaces

